Piecewise-Stationary Combinatorial Bandits

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Motivation

Piecewise-stationary Combinatorial Semi-Bandits

- Problem formulation
- GLR change-point detector
- GLR-CUCB algorithm
- Regret upper bound
- Minimax regret lower bound
- Experimental results

- Problem formulation
- GLRT-CascadeUCB and GLRT-CascadeKL-UCB algorithms
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Motivation

Multi-armed bandit (stochastic) can be defined by K different reward distributions $\{f_1, \ldots, f_k\}$ associated with different arms. **Goal:** Identify the best arm.

Performance metric:

Regret: $\mathcal{R}(T) = T * \mu^* - \mathbb{E}[\sum_{i=1}^{T} \mu_t]$. (convergence rate) Sample Complexity: $n(\epsilon, \delta)$. (final performance) **Application:** Advert placement; Resource allocation; Dynamic pricing.



Many sequential decision making problems have combinatorial nature. Network routing system optimization.



Goal: Minimize the expected total delay.

Top-m arm identification, consider advertise m products to the user.



Goal: Identify *m* most attractive products.

What if the reward distributions are non-stationary?

- Network links might degrade over time.
- User's preference might change.



- Many sequential decision making problems involve combinatorial action, and are in general non-stationary.
- Need a better model for quasi-stationary sequential decision making problems. Existing models are either too optimistic or too pessimistic.
- How good can we perform on this type of problem? Can we achieve optimal performance?

- We study two variants of piecewise-stationary combinatorial bandits, and develop efficient algorithms which achieve nearly order-optimal regret upper bound.
- Key idea of algorithm design is to balance uniform exploration and UCB-type exploration.
- By using randomized hard instance argument, we improve the minimax regret lower bound for piecewise-stationary bandits.

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For $t = 1, \ldots, T$

- Based on the historical data, learner selects a superarm $S_t \in \mathcal{F}$.
- The reward of each base arm contained in the superarm is revealed to the learner $\{R_i(t)|i \in S_t\}$, as well as the reward of the superarm $R_{S_t}(t)$.

Piecewise-stationary CMAB = $(\mathcal{K}, \mathcal{F}, \mathcal{T}, \{f_{k,t}\}_{k \in \mathcal{K}, t \in \mathcal{T}}, r_{\mu_t}(S_t)).$

- \mathcal{K} : set of base arms
- \mathcal{F} : set of super arms
- $\mathcal{T} = \{1, \dots, T\}$: time horizon
- {f_{k,t}}_{k∈K,t∈T}: collection of reward distributions of base arms throughout the time
- $r_{\mu_t}(S_t)$: expected reward function

We assume reward distributions of base arms change in a piecewise manner, let $N = 1 + \sum_{t=1}^{T-1} \mathbb{I}\{\exists k \in \mathcal{K} \text{ s.t. } f_{k,t} \neq f_{k,t+1}\}.$

Goal: identify good super arm at each time step to achieve small regret

Assumption 1. (Monotoncity) Given two arbitrary mean vectors μ and μ' , if $\mu_k \ge \mu'_k$, $\forall k \in \mathcal{K}$, then $r_{\mu}(S) \ge r_{\mu'}(S)$.

Assumption 2. [*L*-Lipschitz] Given two arbitrary mean vectors μ and μ' , there exists an $L < \infty$ such that $|r_{\mu}(S) - r_{\mu'}(S)| \le L ||\mathcal{P}_{S}(\mu - \mu')||_{2}, \forall S \in \mathcal{F}.$

We assume access to an α -approximation oracle $\text{Oracle}_{\alpha}(\mu)$. Given a mean vector μ , $\text{Oracle}_{\alpha}(\mu)$ outputs an α -suboptimal super arm S such that $r_{\mu}(S) \geq \alpha \max_{S \in \mathcal{F}} r_{\mu}(S)$.

Performance metric: Expected α -approximation cumulative regret

$$\mathcal{R}(T) = \mathbb{E}\left[\alpha \sum_{t=1}^{T} \max_{S \in \mathcal{F}} r_{\mu_t}(S) - \sum_{t=1}^{T} r_{\mu_t}(S_t)\right],$$

Consider top-m arm identification. In this case

$$r_{\mu_t}(S_t) = \sum_{i \in S_t} r_i(t)$$

which is the summation of rewards of all base arms contained in the super arm. We can verify that in this case r_{μ_t} is 1-Lipschitiz. The orcale can be realized by any sorting algorithm. In general, $r_{\mu_t}(\cdot)$ can be nonlinear with respect to the rewards of base arms.

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We use generalized likelihood (GLR) change point detector (Besson and Kaufmann, 2019) to monitor the change of base arms' reward distributions.

Algorithm 1: Sub-Bernoulli GLR Change-Point Detector: $GLR(X_1, \dots, X_n; \delta)$

 $\begin{array}{l} \text{Input: observations } X_1, \ldots, X_n \text{ and confidence level } \delta.\\ \text{if } \sup_{s \in [1,n-1]} \quad [s \times \textit{kl}(\hat{\mu}_{1:s}, \hat{\mu}_{1:n}) + (n-s) \times \textit{kl}(\hat{\mu}_{s+1:n}, \hat{\mu}_{1:n})] \geq \beta(n, \delta)\\ \text{then}\\ \mid \quad \text{Return True}\\ \text{end}\\ \text{else}\\ \mid \quad \text{Return False}\\ \text{end}\\ \end{array}$

Advantage: Almost parameter-free.

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GLR-CUCB, runs in three phases:

1. For *p* fraction of the time, the algorithm play a superarm by uniform exploration. For the rest of the time, we play the superarm according to the α -approximation oracle (use UCB indices as input).

2. Once a superarm is played, the algorithm update the UCB indices of all based arm contained in the superarm:

$$\mathsf{UCB}(k) \leftarrow rac{1}{n_\ell} \sum_{n=1}^{n_k} Z_{k,n} + \sqrt{rac{3\log(t-\tau)}{2n_k}}.$$

3. At the end of each round, run GLR change-point detector on all base arms contained in the played superarm. If at least one base arm has changed, restart the UCB indices of all base arms.

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Piecwise-stationary CMAB: suboptimal gaps

The set of bad super arms with respect to the *i*th piecewise-stationary segment is defined as:

$$\mathcal{S}_{B}^{i} = \{ \mathcal{S} | r_{\mu^{i}}(\mathcal{S}) \leq lpha \max_{ ilde{\mathcal{S}} \in \mathcal{F}} r_{\mu^{i}}(ilde{\mathcal{S}}) \}$$

The *suboptimality gaps* in the *i*th stationary segment are:

$$\Delta_{\text{opt}}^{\min,i} = \alpha \max_{\tilde{S} \in \mathcal{F}} r_{\mu^{i}}(\tilde{S}) - \max\{r_{\mu^{i}}(S) | S \in \mathcal{S}_{B}^{i}\},$$

$$\Delta_{\text{opt}}^{\max,i} = \alpha \max_{\tilde{S} \in \mathcal{F}} r_{\mu^{i}}(\tilde{S}) - \min\{r_{\mu^{i}}(S) | S \in \mathcal{S}_{B}^{i}\}.$$

Denote the largest gap at change-point ν_i as

$$\Delta_{\mathsf{change}}^{i} = \max_{k \in \mathcal{K}} \left| \mu_{k}^{i+1} - \mu_{k}^{i} \right|, \forall 1 \leq i \leq N-1.$$

Piecewise-stationary CMAB: regret upper bound

Assumption 3. Define $d_i = d_i(p, \delta) = \left| \{4K/p \left(\Delta_{\text{change}}^i\right)^2\}\beta(T, \delta) + \frac{K}{p} \right|$ and assume $\nu_i - \nu_{i-1} \ge 2 \max\{d_i, d_{i-1}\}, \forall i = 1, \dots, N-1$, where $\nu_N - \nu_{N-1} \ge 2d_{N-1}$. **Remark:** The length of each piecewise-stationary segment is $\Omega(\sqrt{T \log T})$.

Theorem (Zhou et al. 2020)

The expected α -approximation cumulative regret of GLR-CUCB with exploration probability p and confidence level δ satisfies

$$\mathcal{R}(T) \leq \sum_{i=1}^{N} \widetilde{C}_{i} + \Delta_{\text{opt}}^{max} Tp + \sum_{i=1}^{N-1} \Delta_{\text{opt}}^{max,i+1} d_{i} + 3NT \Delta_{\text{opt}}^{max} K\delta,$$

where $\widetilde{C}_{i} = \left(6L^{2} K^{2} \log T / \left(\Delta_{\text{opt}}^{min,i} \right)^{2} + \pi^{2} / 6 + K \right) \Delta_{\text{opt}}^{max,i}.$

Piecewise-stationary CMAB: proof sketch

High-level idea: Recursive regret decomposition. **Step 1:** For stationary case, we have



Step 2: Event decomposition.



Define good events $\{C^{(i)}\}_{i=1}^{N-1}$

$$C^{(i)} = \{ \forall j \leq i, \tau_j \in \{\nu_j + 1, \cdots, \nu_j + d_j \} \}, i \in [N-1].$$

Define $F_i = \{\tau_i > \nu_i\}$ and $D_i = \{\tau_i \le \nu_i + d_i\}$, $\forall 1 \le i \le N - 1$. By definition, $C^{(i)} = F_1 \cap D_1 \cap \cdots \cap F_i \cap D_i$. **Step 3:** Regret decomposition. First decompose with respect to F_1 .

$$\begin{aligned} \mathcal{R}(T) &= \mathbb{E}\left[R(T)\right] = \mathbb{E}\left[R(T)\mathbb{I}\{F_1\}\right] + \mathbb{E}\left[R(T)\mathbb{I}\{\overline{F}_1\}\right] \\ &\leq \mathbb{E}\left[R(T)\mathbb{I}\{F_1\}\right] + T\Delta_{\text{opt}}^{\max}\mathbb{P}(\overline{F}_1) \\ &\leq \mathbb{E}\left[R(\nu_1)\mathbb{I}\{F_1\}\right] + \mathbb{E}\left[R(T-\nu_1)\right] + T\Delta_{\text{opt}}^{\max}K\delta \\ &\leq \widetilde{C}_1 + \Delta_{\text{opt}}^{\max,1}\nu_1p + \mathbb{E}\left[R(T-\nu_1)\right] + T\Delta_{\text{opt}}^{\max}K\delta, \end{aligned}$$

Then decompose $\mathbb{E}[R(T - \nu_1)]$ with respect to $\mathcal{C}^{(1)}$. Repeat the above procedure for $\mathcal{C}^{(2)}, \ldots, \mathcal{C}^{(N-1)}$, we can obtain the desired bound.

Corollary (Zhou et al. 2020)

Let
$$\Delta_{\text{change}}^{min} = \min_{i \in [N-1]} \Delta_{\text{change}}^{i}$$
, we have
(*N* is known) Choosing $\delta = \frac{1}{T}$, $p = \sqrt{\frac{NK \log T}{T}}$, gives
 $\mathcal{R}(T) = \mathcal{O}\left(\frac{NK^2 \log T \Delta_{\text{opt}}^{max}}{(\Delta_{\text{opt}}^{min})^2} + \frac{\sqrt{NKT \log T} \Delta_{\text{opt}}^{max}}{(\Delta_{\text{change}}^{min})^2}\right);$
(*N* is unknown) Choosing $\delta = \frac{1}{T}$, $p = \sqrt{\frac{K \log T}{T}}$, gives
 $\mathcal{R}(T) = \mathcal{O}\left(\frac{NK^2 \log T \Delta_{\text{opt}}^{max}}{(\Delta_{\text{opt}}^{min})^2} + \frac{N\sqrt{KT \log T} \Delta_{\text{opt}}^{max}}{(\Delta_{\text{change}}^{min})^2}\right).$

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Theorem (Zhou et al. 2020)

If $K \ge 3$ and $T \ge M_1 N \frac{(K-1)^2}{K}$, then the worst-case regret for any policy is lower bounded by

 $\mathcal{R}(T) \geq M_2 \sqrt{NKT},$

where
$$M_1 = 1/\log rac{4}{3}$$
, $M_2 = 1/24\sqrt{\log rac{4}{3}}$.

Proof sketch:

1. Construct randomized hard instance. $\mu_i^{i^*} \sim \text{Bern}(\frac{1}{2} + \epsilon)$, $\forall k \in \mathcal{K} \setminus i^*, \mu_i^k \sim \text{Bern}(\frac{1}{2})$. $(i+1)^* | i^* \sim \text{uniform}(\mathcal{K} \setminus i^*)$

2. By change of measure technique, one can show that this ensemble of hard instances incurs regret at least $\Omega(\sqrt{NKT})$.

3. There exists at least one instance incurs regret on the order of $\Omega(\sqrt{NTK})$.

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Piecewise-statinoary CMAB: experimental results



Figure: Experiment on synthetic dataset. Left: reward distribution of base arms. Right: expected accumulative regret.

Piecewise-statinoary CMAB: experimental results



Figure: Yahoo! 1 experiment. Left: reward distribution of base arms. Right: expected accumulative regret.

Piecewise-statinoary CMAB: experimental results



Figure: Yahoo! 2 experiment. Left: reward distribution of base arms. Right: expected accumulative regret.

	CUCB	CTS	Hybrid	DUCB	GLR-CUCB
Synthetic Dataset	241.08	351.12	278.82	14.08	37.96
Yahoo! Experiment 1	510.20	513.41	826.01	25.44	62.76
Yahoo! Experiment 2	563.54	562.24	1189.17	158.27	517.13

Table: Standard deviations of all algorithms for experiments on synthetic and Yahoo! datasets

	LR-GLR-CUCB	MUCB	Oracle-CUCB
Synthetic Dataset	73.30	171.45	25.54
Yahoo! Experiment 1	63.37	202.44	35.08
Yahoo! Experiment 2	496.73	1427.28	160.76

Table: Standard deviations of all algorithms for experiments on synthetic and Yahoo! datasets

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For $t = 1, \ldots, T$

- Given historical data, the learner selects *K* out of *L* items to recommend to the user.
- The learner observes a partial feedback of his decision, $\arg\min_k 1 \le k \le K : Z_{a_{k,t}}, t = 1$, the first item clicked by the user/no click.

Piecewise-stationary cascading bandit: problem formulation

Piecewise-stationary cascasding bandit (CB)= $(\mathcal{L}, \mathcal{T}, \{f_{\ell,t}\}_{\ell \in \mathcal{L}, t \in \mathcal{T}}, K)$.

- \mathcal{L} : Ground set containing L items (e.g., web pages or advertisements).
- $\mathcal{T} = \{1, \ldots, T\}$: Set of time steps.
- $\{f_{\ell,t}\}_{\ell \in \mathcal{L}, t \in \mathcal{T}}$: Pmfs of items in \mathcal{L} at all time steps.
- K: Number of items recommended by the learner to the user.

Learner receives partial feedback at time t, given by

$$F_t = \begin{cases} \emptyset, & \text{if no click,} \\ \arg\min_k \{1 \le k \le K : Z_{\mathbf{a}_{k,t},t} = 1\}, & \text{otherwise.} \end{cases}$$

Goal of learner: Identify top-K items with highest clicked probabilities.

$$\mathcal{R}(T) = \mathbb{E}\left[\sum_{t=1}^{T} R\left(\mathcal{A}_{t}, \mathbf{w}_{t}, \mathbf{Z}_{t}\right)\right],$$

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GLRT-CascadeUCB and GLRT-CascadeKL-UCB algorithm run in three phases:

1. For p fraction of the time, the algorithm select K items by uniform sampling. For the rest of the time, the algorithm select K items with highest UCB/KL-UCB indices.

2. Update the statistics of *K* selected items.

$$\begin{aligned} \mathsf{UCB}(\ell) &= \hat{\mathbf{w}}(\ell) + \sqrt{\frac{3\log(t-\tau)}{2n_{\ell}}}, \\ \mathsf{UCB}_{\mathsf{KL}}(\ell) &= \max\{q \in [\hat{\mathbf{w}}(\ell), 1] : n_{\ell} \times \mathsf{KL}(\hat{\mathbf{w}}(\ell), q) \leq g(t-\tau)\}. \end{aligned}$$

3. At the end of each round, run GLR change-point detector on selected items at this round. If at least one item's click probability has changed, restart the UCB indices/KL-UCB indices of all items.

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Theorem (Wang et al. 2019)

GLRT-CascadeUCB guarantees

$$\mathcal{R}(T) \leq \sum_{\substack{i=1\\(a)}}^{N} \widetilde{C}_{i} + \underbrace{Tp}_{(b)} + \underbrace{\sum_{\substack{i=1\\(c)}}^{N-1} d_{i} + \underbrace{3NTL\delta}_{(d)},$$
where $\widetilde{C}_{i} = \sum_{\ell=K+1}^{L} \frac{12}{\Delta_{s_{i}(\ell),s_{i}(K)}^{i}} \log T + \frac{\pi^{2}}{3}L.$

Theorem (Wang et al. 2019)

GLRT-CascadeKL-UCB guarantees

$$\mathcal{R}(T) \leq \underbrace{\mathcal{T}(N-1)(L+1)\delta}_{(a)} + \underbrace{\mathcal{T}p}_{(b)} + \underbrace{\sum_{i=1}^{N-1} d_i}_{(c)} + \underbrace{\mathcal{N}K \log \log T + \sum_{i=0}^{N-1} \widetilde{D}_i}_{(d)},$$

where \widetilde{D}_i is a term depending on log T and the suboptimal gaps.

Corollary (Wang et al. 2019)

The regret of GLRT-CascadeUCB is established by choosing $\delta = \frac{1}{T}$ and $p = \sqrt{\frac{NL \log T}{T}}$:

$$\mathcal{R}(T) = \mathcal{O}\left(\frac{N(L-K)\log T}{\Delta_{\text{opt}}^{\min}} + \frac{\sqrt{NLT\log T}}{\left(\Delta_{\text{change}}^{\min}\right)^2}\right).$$
 (1)

Choosing the same δ and p , <code>GLRT-CascadeKL-UCB</code> has same order of regret upper bound as (1).

remark: the order of the regret upper bound is the same as GLR-CUCB, which implies that the dominant factor is the change in distribution.

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Piecewise-stationary cascading bandit: minimax regret lower bound

Theorem (Wang et al. 2019)

If $L \ge 3$ and $T \ge MN\frac{(L-1)^2}{L}$, then for any policy, the worst-case regret is at least $\Omega(\sqrt{NLT})$, where $M = 1/\log \frac{4}{3}$, and $\Omega(\cdot)$ notation hides a constant factor that is independent of N, L, and T.

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Piecewise-stationary cascading bandit: experimental results



Figure: Synthetic experiment. Left: reward distributions. Right: cumulative regret.

Piecewise-stationary cascading bandit: experimental results



Figure: Experiment on Yahoo! dataset. Left: reward distributions. Right: cumulative regret.

Table: Means and standard deviations of the *T*-step regrets.

	CascadeUCB1	CascadeKL-UCB	CascadeDUCB
Synthetic Dataset	1069.77 ± 87.09	1053.25 ± 111.67	1180.30 ± 20.22
Yahoo! Experiment	2349.29 ± 312.71	2820.16 ± 256.74	3226.97 ± 39.37
	CascadeSWUCB	GLRT-CascadeUCB	GLRT-CascadeKL-UCB
Synthetic Dataset	664.84 ± 29.81	$\textbf{527.93} \pm \textbf{25.20}$	$\textbf{440.93} \pm \textbf{45.54}$
Yahoo! Experiment	1519.56 ± 52.23	1235.21 ± 54.59	$\textbf{856.77} \pm \textbf{67.16}$
	Oracle-CascadeUCB1	Oracle-CascadeKL-UC	
Synthetic Dataset	472.25 ± 17.65	353.86 ± 19.59	
Yahoo! Experiment	1230.17 ± 45.24	808.84 ± 47.97	

- We develop the first efficient algorithm for piecewise-stationary combinatorial semi-bandits, GLR-CUCB, which achieves $\mathcal{O}(\sqrt{NKT \log T})$ regret.
- We improve minimax regret lower bound $(\Omega(\sqrt{NKT}))$ for piecewise-stationary combinatorial semi-bandits, which indicates GLR-CUCB is nearly order-optimal within poly-logarithm factors.
- We develop better algorithms for piecewise-stationary cascading bandits and tighten the minimax regret lower bound.
- Future work includes design time-unaware algorithms for piecewise-stationary bandits, incorporate contextual information, etc.

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Q & A

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